

Amendments to the Claims

This listing of claims will replace all prior versions, and listings, of claims in the application:

Listing of Claims:

Claim 1. (currently amended) A computer system providing an object-based virtual machine environment for running successive applications, said computer system ~~including~~ comprising:

storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next, ~~said system including~~ :

a card table comprising multiple cards, each corresponding to a region of said storage, each card in the card table being set to null when the first heap is reset between successive applications;

~~means~~ logic for marking a card whenever an object in its corresponding storage region is updated; and

~~means~~ logic for detecting possible references from the second heap to the first heap at reset by scanning the cards in the card table corresponding to the second heap, and detecting any cards which have been marked.

Claim 2. (currently amended) The computer system of claim 1, further comprising: ~~means~~ logic for locating, for each marked card, any objects in the corresponding region of storage; and ~~means~~ logic for identifying any references to the first heap in the located objects.

Claim 3. (currently amended) The computer system of claim 2, further comprising: ~~means~~ logic responsive to the identification of references to the first heap for performing the mark phase of a garbage collection to determine live objects in at least the second heap; ~~means~~ logic for detecting whether any objects in the second heap having references to the first heap have been marked as live; and ~~means~~ logic responsive to a detection of any such objects for returning an error

condition to prevent reset for another application.

Claim 4. (currently amended) The computer system of claim 3, further comprising ~~means~~ logic for invalidating the card table if a compact operation has been performed on the second heap since the last reset, wherein said means for performing the mark phase is also responsive to invalidation of the card table.

Claim 5. (previously presented) The computer system of claim 1, wherein an object is only considered as within the region of storage corresponding to a card if a predetermined part of the object is in that region.

Claim 6. (previously presented) The computer system of claim 1, wherein the region of memory corresponding to a card comprises between 256 and 2048 bytes.

Claim 7. (currently amended) The computer system of claim 1, further comprising: ~~means~~ logic for detecting references or possible references to the first heap from a set of predetermined locations; and ~~means~~ logic responsive to the detection of any such references or possible references for returning an error condition to prevent reset for another application.

Claim 8. (currently amended) The computer system of claim 7, wherein ~~said~~ the set of predetermined locations includes the stacks and registers.

Claim 9. (currently amended) The computer system of claim 1, further comprising: ~~means~~ logic for detecting any objects on the first heap which are reachable from virtual machine system class objects; and ~~means~~ logic for promoting any such detected objects to the second heap.

Claim 10. (currently amended) A computer system providing an object-based virtual machine environment for running successive applications, said computer system ~~including~~ comprising: storage, at least a portion of which is logically divided into two or more heaps in which

objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next; ~~said system including:~~

means for identifying any objects on the first heap which have a finalization method; and
means for running the finalization methods of any identified objects on the main thread prior to reset of the first heap.

Claim 11. (currently amended) The computer system of claim 10, further comprising ~~means~~ logic responsive to running ~~said~~ the finalization methods for checking that they have not performed any operations which would prevent reset of the first heap.

Claim 12. (currently amended) A method of operating a computer system providing an object-based virtual machine environment for running successive applications, ~~said~~ the computer system including storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next, said method including the steps of:

providing a card table comprising multiple cards, each corresponding to a region of ~~said~~ the storage, each card in the card table being set to null when the first heap is reset between successive applications;

marking a card whenever an object in its corresponding storage region is updated; and
detecting possible references from the second heap to the first heap at reset by scanning the cards in the card table corresponding to the second heap, and detecting any cards which have been marked.

Claim 13. (previously presented) The method of claim 12, further comprising: locating, for each marked card, any objects in the corresponding region of storage; and identifying any references to the first heap in the located objects.

Claim 14. (currently amended) The method of claim 13, further comprising: responsive to the identification of references to the first heap, performing the mark phase of a garbage collection to

determine live objects in at least the second heap; detecting whether any objects in the second heap having references to the first heap have been marked as live; and responsive to a detection of any such objects, returning an error condition to prevent reset for another application.

Claim 15. (previously presented) The method of claim 14, further comprising the step of invalidating the card table if a compact operation has been performed on the second heap since the last reset, wherein said step of performing the mark phase is also performed in response to invalidation of the card table.

Claim 16. (previously presented) The method of claim 12, wherein an object is only considered as within the region of storage corresponding to a card if a predetermined part of the object is in that region.

Claim 17. (previously presented) The method of claim 12, wherein the region of memory corresponding to a card comprises between 256 and 2048 bytes.

Claim 18. (previously presented) The method of claim 12, further comprising: detecting references or possible references to the first heap from a set of predetermined locations; and responsive to the detection of any such references or possible references, returning an error condition to prevent reset for another application.

Claim 19. (previously presented) The method of claim 18, wherein said set of predetermined locations includes the stacks and registers.

Claim 20. (previously presented) The method of claim 12, further comprising: detecting any objects on the first heap which are reachable from virtual machine system class objects; and promoting any such detected objects to the second heap.

Claim 21. (previously presented) A method of operating a computer system providing an object-

based virtual machine environment for running successive applications, said computer system including storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next, said method including the steps of: identifying any objects on the first heap which have a finalization method; and running the finalization methods of any identified objects on the main thread prior to reset of the first heap.

Claim 22. (previously presented) The method of claim 21, further comprising the step, responsive to running said finalization methods, of checking that they have not performed any operations which would prevent reset of the first heap.

Claim 23. (currently amended) A computer program product for operating a computer system providing an object-based virtual machine environment for running successive applications, ~~said~~ the computer system including storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next, said computer program product comprising machine-readable program instructions recorded on a storage medium, ~~said~~ the instructions when loaded into the computer system causing it to perform the steps of: providing a card table comprising multiple cards, each corresponding to a region of ~~said~~ the storage, each card in the card table being set to null when the first heap is reset between successive applications; marking a card whenever an object in its corresponding storage region is updated; and detecting possible references from the second heap to the first heap at reset by scanning the cards in the card table corresponding to the second heap, and detecting any cards which have been marked.

Claim 24. (previously presented) The computer program product of claim 23, wherein said program instructions further cause the computer to perform the steps of: locating, for each marked card, any objects in the corresponding region of storage; and identifying any references to the first heap in the located objects.

Claim 25. (previously presented) The computer program product of claim 24, wherein said program instructions further cause the computer to perform the steps of: responsive to the identification of references to the first heap, performing the mark phase of a garbage collection to determine live objects in at least the second heap; detecting whether any objects in the second heap having references to the first heap have been marked as live; and responsive to a detection of any such objects, returning an error condition to prevent reset for another application.

Claim 26. (previously presented) The computer program product of claim 25, wherein said program instructions further cause the computer to perform the step of invalidating the card table if a compact operation has been performed on the second heap since the last reset, wherein said step of performing the mark phase is also performed in response to invalidation of the card table.

Claim 27. (previously presented) The computer program product of claim 23, wherein an object is only considered as within the region of storage corresponding to a card if a predetermined part of the object is in that region.

Claim 28. (previously presented) The computer program product of claim 23, wherein the region of memory corresponding to a card comprises between 256 and 2048 bytes.

Claim 29. (previously presented) The computer program product of claim 23, wherein said program instructions further cause the computer to perform the steps of: detecting references or possible references to the first heap from a set of predetermined locations; and responsive to the detection of any such references or possible references, returning an error condition to prevent reset for another application.

Claim 30. (previously presented) The computer program product of claim 29, wherein said set of predetermined locations includes the stacks and registers.

Claim 31. (previously presented) The computer program product of claim 23, wherein said program instructions further cause the computer to perform the steps of: detecting any objects on the first heap which are reachable from virtual machine system class objects; and promoting any such detected objects to the second heap.

Claim 32. (previously presented) A computer program product for operating a computer system providing an object-based virtual machine environment for running successive applications, said computer system including storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next, said computer program product comprising machine-readable program instructions recorded on a storage medium, said instructions when loaded into the computer system causing it to perform the steps of: identifying any objects on the first heap which have a finalization method; and running the finalization methods of any identified objects on the main thread prior to reset of the first heap.

Claim 33. (previously presented) The computer program product of claim 21, wherein said program instructions further cause the computer to perform the step responsive to running said finalization methods, checking that they have not performed any operations which would prevent reset of the first heap.

Claim 34. (new) A computer system providing an object-based virtual machine environment for running successive applications, said computer system comprising:

storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next;

a card table comprising multiple cards, each corresponding to a region of said storage, each card in the card table being set to null when the first heap is reset between successive applications;

logic for marking a card whenever an object in its corresponding storage region is updated;

logic for detecting possible references from the second heap to the first heap at reset by scanning the cards in the card table corresponding to the second heap, and detecting any cards which have been marked;

logic for locating, for each marked card, any objects in the corresponding region of storage; and logic for identifying any references to the first heap in the located objects;

logic responsive to the identification of references to the first heap for performing the mark phase of a garbage collection to determine live objects in at least the second heap; logic for detecting whether any objects in the second heap having references to the first heap have been marked as live; and logic responsive to a detection of any such objects for returning an error condition to prevent reset for another application;

logic for invalidating the card table if a compact operation has been performed on the second heap since the last reset;

logic for detecting references or possible references to the first heap from a set of predetermined locations; and logic responsive to the detection of any such references or possible references for returning an error condition to prevent reset for another application;

storage, at least a portion of which is logically divided into two or more heaps in which objects can be stored, wherein a first heap is reset between successive applications, and a second heap persists from one application to the next;

logic for identifying any objects on the first heap which have a finalization method; and

logic for running the finalization methods of any identified objects on the main thread prior to reset of the first heap;

wherein an object is only considered as within the region of storage corresponding to a card if a predetermined part of the object is in that region; and wherein the region of memory corresponding to a card comprises between 256 and 2048 bytes; and wherein the logic for performing the mark phase is also responsive to invalidation of the card table.